

When you're connected by zoom:

- ▶ Use a laptop or desktop with a large screen so you can read these words clearly.
- ▶ To conserve bandwidth, please turn off your camera.
- ▶ Please mute your microphone unless I call on you.
- ▶ Please have the chat window open to ask questions.
- ▶ Take-home final due in 1 week. (But all deadlines are elastic.)
- ▶ Today's DJ: Trent.

Take-home final out by Wed.

# Kurosh Subgroup Theorem

## Theorem (Kurosh)

*Every subgroup of a free group is free.*

For simplicity, we discuss and prove the finite index case, though everything works in general.

# The fundamental group of a basepointed $X$ -labelled graph

Let  $\Gamma$  be an  $X$ -labelled graph with basepoint  $v_0$ .

## Definition

The **fundamental group** of  $\Gamma$ , written  $\pi_1(\Gamma, v_0)$ , is:

- ▶ Elements: Reduced loops starting at  $v_0$ .
- ▶ Operation: Concatenation.
- ▶ Identity: Empty word 1.
- ▶ Inverse of a loop: The same loop, traversed backwards.

**Observation:** If we change the labels on  $\Gamma$ ,  $\pi_1(\Gamma, v_0)$  doesn't change; we just have different names for the same paths.

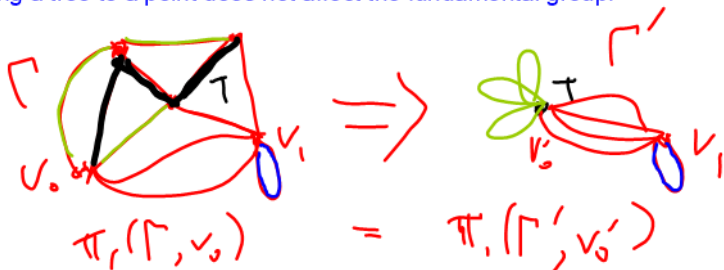
$$\pi_1(\Gamma, v_0) = \pi_1(\Gamma, v_0)$$

# The Retract Lemma

## Theorem

Let  $\Gamma$  be a basepointed  $X$ -labelled graph, let  $T$  be a tree contained in  $\Gamma$ , and let  $\Gamma'$  be the basepointed labelled graph obtained by replacing  $T$  with a single vertex and giving each edge a different label in some alphabet  $Y$ . Then  $\pi_1(\Gamma, v_0) = \pi_1(\Gamma', v_0')$ .

Collapsing a tree to a point does not affect the fundamental group.



# Proof of the Retract Lemma

and concatenations map to concatenations

First: Relabel  $\Gamma$  so that every edge has a different label.

Loops map to loops, so enough to show that:

- ▶ (Surjective) Every loop in  $\Gamma'$  is the image of some loop in  $\Gamma$ .
- ▶ (Injective) Nontrivial reduced loops map to nontrivial reduced loops.

in  $\Gamma'$

in  $\Gamma$

— outside  $T$   
— inside  $T$   
—  $T$ , not  $T$

Surjectivity: Suppose we have a loop in  $\Gamma'$ . Illustrate using example.



pull back

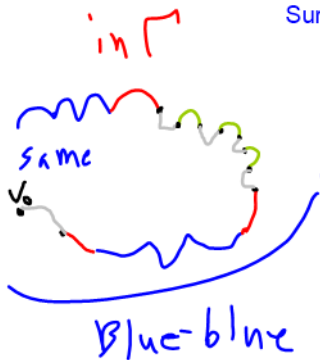


Loop in  $\Gamma'$



Red-green-red paths may pull back to disconnected pieces; similarly, a path that ends in a red edge in  $\Gamma'$  might not return to basepoint when pulled back.

But remember, any two points in a tree  $T$  are connected by a unique path within  $T$ . So we can fill all those gaps with paths from within  $T$  to get a loop starting and ending at  $v_0$ :



Surjectivity follows.

Injectivity: Suppose we have a reduced loop in  $\Gamma$ . How could this map to a non-reduced loop in  $\Gamma'$ ? Look at all possible cases of consecutive edges in the loop in  $\Gamma'$ :

Blue-blue



These edges inject



, so can't cancel

Blue-red



different color edges so different labels

Same argument works for green-red.

Green-green



from

(this is mapped into Gamma')



If this happened w/reduced black path, then we would lose injectivity.

Fortunately, there is a unique reduced path between any two points in a tree, including the case where the two points are the same. So the picture in Gamma is actually one where the black path is trivial:



So actually, we didn't have a reduced path in Gamma; contradiction. Same argument holds for red-red cancellation.

One case left: What if we have a reduced loop that's entirely black? Well, there's only one such path, namely, the trivial path. So the only element of the kernel of our homomorphism is the trivial path.

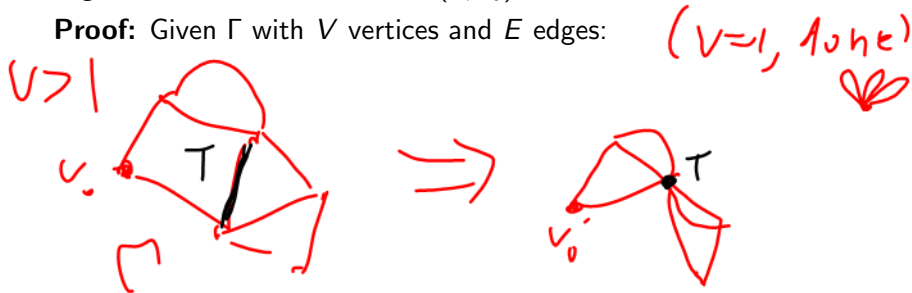


# Fundamental groups of graphs are free

## Corollary

Let  $\Gamma$  be a finite basepointed  $X$ -labelled graph with basepoint  $v_0$ ,  $E$  edges, and  $V$  vertices. Then  $\pi_1(\Gamma, v_0) = F_n$ , where  $n = E - V + 1$ .

**Proof:** Given  $\Gamma$  with  $V$  vertices and  $E$  edges:



Lose one edge, one vertex, so  $E-V+1$  remains constant. (Actually  $V-E$  is Euler)  
Result follows by induction on  $V$ .

Same proof works for infinite graphs, but you have to prove things about spanning trees of infinite graphs.



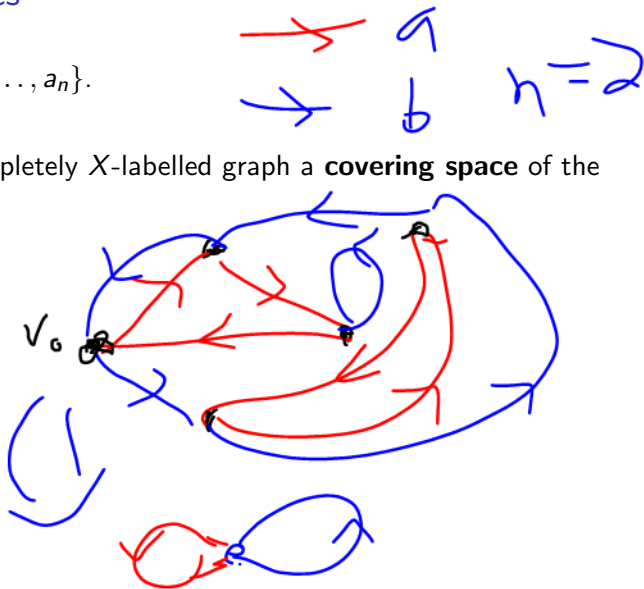
# Covering spaces

Let  $X = \{a_1, \dots, a_n\}$ .

## Definition

We call a completely  $X$ -labelled graph a **covering space** of the  $n$ -leaved rose:

graph map  
induced by  
colors & arrows

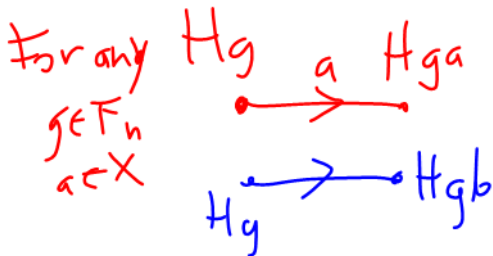


# Coset diagrams

## Definition

$$n=2$$

For  $H \leq F_n$ , we define the (right) **coset diagram** of  $H$  to be the basepointed  $X$ -labelled graph given by:



Gives completely  $X$ -labelled graph whose vertices are the right cosets of  $H$ .

Basept:  $H$

# Main Theorem of covering spaces

Let  $X = \{a_1, \dots, a_n\}$ .

## Theorem

There exists a bijection between subgroups of  $F_n$  and basepointed covering spaces of the  $n$ -leaved rose, given by:

- ▶ **Subgroup to covering space:** Map  $H$  to its coset diagram. *basepointed*
- ▶ **Covering space to subgroup:** Map  $(\Gamma, v_0)$  to  $\pi_1(\Gamma, v_0)$ .

*Proof:*

$H \rightarrow$  coset diag  
basept  $H$

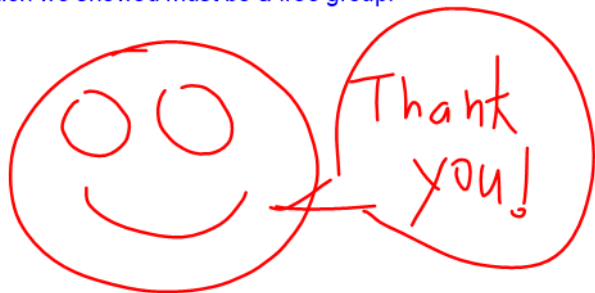
Check: These maps are inverses.

$\pi_1(\Gamma, v_0) \leftarrow (\Gamma, v_0)$

# Proof of the Kurosh Subgroup Theorem

$H$  subgroup of finite index in  $F_n$ .

By Correspondence of Covering Spaces,  $H$  is isomorphic to fundamental group of a graph, which we showed must be a free group.



# Residual finiteness of free groups

Here's a slightly fancier result about free groups.

## Theorem

*Given  $1 \neq w \in F_n$ , there exists a finite quotient  $\rho : F_n \rightarrow G$  such that  $\rho(w) \neq 1$ .*

We say that free groups are **residually finite**: Given a nontrivial element  $w$  of a free group  $F_n$ , there exists a finite quotient of  $F_n$  in which some nontrivial “residue” survives.

# Reducing the problem

Enough to show that:

## Theorem

*Given  $1 \neq w \in F_n$ , there exists a subgroup  $H$  of finite index in  $F_n$  such that  $w \notin H$ .*

For in that case, let  $N$  be the intersection of the finitely many conjugates of  $H$ . This is a subgroup of finite index in  $F_n$ , and  $\rho : F_n \rightarrow F_n/N$  is the desired finite quotient.

## Proof by example

In  $F_2 = \langle a, b \rangle$ , take  $w =$