

Math 127, Mon Mar 22

- ▶ Use a laptop or desktop with a large screen so you can read these words clearly.
- ▶ In general, please turn off your camera and mute yourself.
- ▶ Exception: When we do groupwork, please turn both your camera and mic on. (Groupwork will not be recorded.)
- ▶ Please always have the chat window open to ask questions.
- ▶ Reading for today: 7.1–7.3.
- ▶ PS06 due tonight, late deadline Fri Mar 26.
- ▶ **Exam 2 Wed Mar 24**, on 3.5–3.6, 4.2–4.3, 5.3–5.6, and 6.1–6.4 (PS04–06). Review session at 3pm (recorded to YouTube).

Questions on Ch. 6 and PS06?

Make sure you know the following chain of definitions:

- * To say H is parity check matrix of a binary linear code C means that $C = \text{Null}(H)$.
- * $\text{Null}(H)$ is the set of all x such that $Hx = 0$. In other words, in the context of binary linear codes, we think of H as the matrix of a system of linear equations (over F_2), and the code C is the solution set for that system of linear equations. If we write out a basis for $\text{Null}(H)$ as the columns of a matrix G , that matrix G is a generator matrix for C .

Again, know your definitions:

- * binary linear code
- * parity check matrix
- * generator matrix
- * $\text{Null}(H)$

For example, 6.3.2(b): Need to list all of the codewords in a code C given by a parity check matrix H .

To do that:

- * Think of H as the matrix of a system of linear equations
- * Solve H and get a basis for $\text{Null}(H)=C$
- * Use that basis to list all possible vectors in C .

Q: Do you treat parity check bits differently from data bits?

A: No.



By isolating just the transmission part, we can get a better understanding of what is possible in an error-correcting code.

Ideals

\mathbb{Z} or $F[x]$

Maybe the most important definition in ring theory:

Definition

Let R be a (commutative) ring. An **ideal** of R is $I \subseteq R$ s.t.:

1. (Zero) The zero element of R is contained in I .
2. (Closed under addition) If $x, y \in I$, then $x + y \in I$.
3. (Closed under R -multiplication) If $x \in I$ and $r \in R$, then $rx \in I$.

Exs:

Compare: Definition of subspace/subspace test

For a ring R :

- ▶ The set $\{0\}$ is an ideal of R called the **zero ideal**.
- ▶ R is an ideal of itself.

\mathbb{Z}

More interesting examples

integer multiples of 3

$$\text{Let } R = \mathbf{Z}, I = \{3n \mid n \in \mathbf{Z}\} = \{\dots, -9, -6, -3, 0, 3, 6, 9, \dots\} = (3)$$

$$0 \in I \quad \checkmark$$

$$\text{If } x, y \in I, x = 3n, y = 3k \Rightarrow x + y = 3(n+k) \in I \quad \checkmark$$

$$\text{If } x \in I, r \in \mathbf{Z}, x = 3n, rx = 3nr \in I. \quad \checkmark$$

$$\text{Let } R = \mathbf{F}_2[x], I = \{p(x)(x^2 + x + 1) \mid p(x) \in \mathbf{F}_2[x]\}.$$

$$= (x^2 + x + 1)$$

So I is an ideal of the ring \mathbf{Z} .

I.e. $I =$ all poly mults of $(x^2 + x + 1)$

$$= \{0, x^2 + x + 1, x^3 + x^2 + x, (x+1)(x^2 + x + 1) = x^3 + 1, \dots\}$$

Check I is ideal of $\mathbb{F}_2[x]$:

$$0 \in I \quad \checkmark$$

$$\textcircled{A} \quad y, z \in I \Rightarrow \begin{aligned} y &= p(x)(x^2+x+1) \\ z &= q(x)(x^2+x+1) \end{aligned}$$

$$\Rightarrow y+z = (p(x)+q(x))(x^2+x+1) \in I \quad \checkmark$$

$$\textcircled{A} \quad y \in I, r(x) \in \mathbb{F}_2[x]$$

$$\Rightarrow y = p(x)(x^2+x+1)$$

$$ry = \underbrace{(r(x)p(x))}_{\in \mathbb{F}_2[x]}(x^2+x+1) \in I \quad \checkmark$$

Classes of examples

R a ring.

$$R = \mathbb{Z}, F[x]$$

- ▶ For fixed $a \in R$, the set

$$(a) = \{ra \mid r \in R\} = \{\text{all } R\text{-multiples of } a\}$$

is called the **principal ideal generated by a** .

- ▶ For fixed $a, b \in R$, the set

$$(a, b) = \{ra + sb \mid r, s \in R\}$$

is called the **ideal generated by a and b** .

- ▶ For F a field and $a \in F$, the set

$$I_a = \{f(x) \in F[x] \mid f(a) = 0\}$$

is an ideal of $F[x]$.

Quotient rings

One of the most important uses of ideals is to mod out by them.
Specifically:

How can we make sense of $F[x]/(m(x))$ the same way we made sense of $\mathbf{Z}/(m)$?

(And to be honest, we never really addressed all of the details of making sure that $\mathbf{Z}/(m)$ works, so we'll do that too.)

Cosets

Cosets are a fancy way of defining modular equivalence!

Let R be a ring, and let I be an ideal of R . For $r \in R$, we define the **additive coset** $r + I$ to be

$$r + I = \{r + a \mid a \in I\}.$$

If the context is clear, instead of saying “additive coset”, we just say **coset**.

Example: What are the cosets of (3) in \mathbb{Z} ?

$0 + I = \{\dots, -6, -3, 0, 3, 6, 9, \dots\}$ $\left. \begin{array}{l} 0 \\ \text{mod} \\ 3 \end{array} \right\}$

$4 + I = \{\dots, -5, -2, 1, 4, 7, 10, 13, \dots\}$ $\left. \begin{array}{l} \text{mod} \\ 3 \end{array} \right\}$

$\quad = 1 + I = 10 + I$

$8 + I = \{\dots, 2, 5, 8, 11, 14, \dots\}$ $\left. \begin{array}{l} 2 \\ \text{mod} \\ 3 \end{array} \right\}$

$\quad = 2 + I$

Note: A red "= I" is written above the example text.

When are elements in the same coset?

Theorem

R be a ring, I be an ideal of R , $r, s \in R$. TFAE:

1. $r + I = s + I$ (i.e., the cosets $r + I$ and $s + I$ are the same set).
2. $r \in s + I$.
3. $r - s \in I$.
4. $r = s + a$ for some $a \in I$.

Proof: (1) \Rightarrow (2) \Rightarrow (3) \Rightarrow (4); (4) \Rightarrow (1) on HW.

(A) (2) $r \in s + I$

So $r = s + a$ for some $a \in I$

$\Rightarrow r - s = a$ " " " "

(C) (3) $r - s \in I$.

$$\textcircled{A} \quad (3) \quad r - s \in I$$

So $r - s = a$ for some $a \in I$

$$\Rightarrow r = s + a \quad " \quad " \quad " .$$

$$\textcircled{C} \quad (4) \quad r = s + a \text{ for some } a \in I$$

$$\text{I.e. } r + I = s + I$$

$$\Leftrightarrow r = s + a, a \in I.$$


Definition

To say that r is a **representative of the coset** $s + I$ means that $r \in s + I$.

Definition of quotient ring

Let R be a ring and let I be an ideal of R . We define the **quotient ring** R/I as follows.

- ▶ **Set:** The elements of R/I are the cosets of I in R . Note that if r and s represent the same coset of I , then the cosets $r + I$ and $s + I$ are actually the same element of R/I , since $r + I = s + I$.
- ▶ **Addition:** For $r + I, s + I \in R/I$, we define the sum

$$(r + I) + (s + I) = (r + s) + I.$$

- ▶ **Multiplication:** For $r + I, s + I \in R/I$, we define the product

$$(r + I)(s + I) = rs + I.$$

The zero element of R/I is $0 + I = I$, and the one element is $1 + I$.

Example: $\mathbf{Z}/(3)$

Let $R = \mathbf{Z}$, $I = (3)$ (the multiples of 3 in \mathbf{Z}).

Elements:

$$0+I, 1+I, 2+I$$

Multiplication table:

	$0+I$	$1+I$	$2+I$
$0+I$	$0+I$	$0+I$	$0+I$
$1+I$	$0+I$	$1+I$	$2+I$
$2+I$	$0+I$	$2+I$	$1+I$

$$(0+I)(r+I) = 0+I$$

$$(1+I)(r+I) = (1 \cdot r) + I$$

$$= r+I$$

$$(2+I)(2+I) = 4+I = 1+I$$

So this really is the same $\mathbf{Z}/(3)$ we've been talking about all along.

The potential problem with quotients

When we define:

$$(r + I) + (s + I) = (r + s) + I \quad (r + I)(s + I) = (rs) + I$$

Could it be the case that you get a different answer if you use different representatives for the cosets $r + I$ and $s + I$?

Theorem: No, everything works fine. (I.e., the operations in a quotient ring are well-defined and don't depend on our choice of coset representative.)

~~**Proof:** (multiplication part)~~

(uses $v + I = s + I$)

$(\Leftrightarrow v = s + \alpha \text{ for } \alpha \in I)$

Review/revision: Computation in $\mathbf{Z}/(m)$

Let $I = (m)$ (the integer multiples of m). Working mod I , we have:

- ▶ **Elements:** The cosets of I in \mathbf{Z} , which we can write as $0 + I, 1 + I, \dots, (m - 1) + I$, or $\{0, \dots, m - 1\}$ for short, since division by m gives remainders between 0 and $m - 1$.
- ▶ **Operations:** Addition and multiplication are computed in \mathbf{Z} and then reduced mod I . I.e., you use division by m with remainder to choose a **reduced representative** for the final answer.

Example:

Computation in $F[x]/(m(x))$, version 1

F a field, $m(x) \in F[x]$ ($\deg m > 0$), $I = (m(x))$ (the polynomial multiples of $m(x)$). Working mod I , we have:

- ▶ **Elements:** The cosets of I in $F[x]$, which we can write as $r(x) + I$ where $\deg r(x) < \deg m(x)$, since division by $m(x)$ gives remainders of degree $< \deg m(x)$.
- ▶ **Operations:** Addition and multiplication are computed in $F[x]$ and then reduced mod I . I.e., you use division by $m(x)$ with remainder to choose a **reduced representative** for the final answer.

Example:

Computation in $F[x]/(m(x))$, version 2

F a field, $m(x) \in F[x]$ ($\deg m > 0$), $I = (m(x))$ (the polynomial multiples of $m(x)$). Abbreviate $\alpha = x + I$. Working mod I , we have:

- ▶ **Elements:** The cosets of I in $F[x]$, which we can write as $r(\alpha)$ where $\deg r < \deg m$, since setting $m(\alpha) = 0$ allows you to reduce any polynomial of degree $\geq \deg m$.
- ▶ **Operations:** Addition and multiplication are computed in polynomials in α and then reduced. I.e., you use the relation $m(\alpha) = 0$ to choose a **reduced representative** for the final answer.

Example: